Learning-based Analysis and Control of Safety-Critical Systems

Enrique Mallada



Autonomous Energy Systems Workshop NREL

July 14, 2022

A World of Success Stories

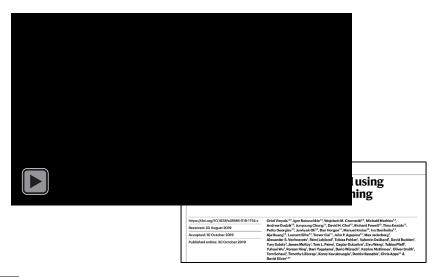
2017 Google DeepMind's DQN

2017 AlphaZero – Chess, Shogi, Go



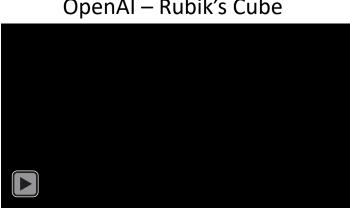
Boston Dynamics





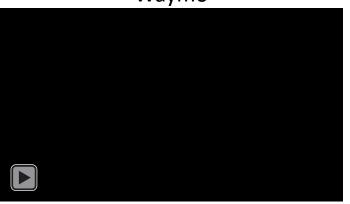


ETTER





Waymo



Reality Kicks In

Angry Residents, Abrupt Stops: Waymo Vehicles Are Still Causing Problems in Arizona



DeepMind's Losses and the Future of Artificial Intelligence

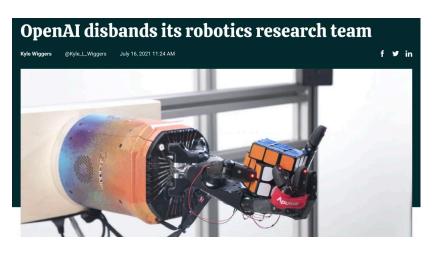
Alphabet's DeepMind unit, conqueror of Go and other games, is losing lots of money. Continued deficits could imperil investments in Al.

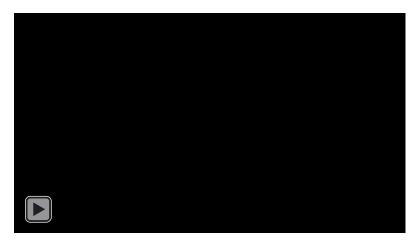
AARIAN MARSHALL BUSINESS 12.87.2020 04:06 PM

Uber Gives Up on the Self-Driving Dream

RAY STERN | MARCH 31, 2021 | 8:26AM

The ride-hail giant invested more than \$1 billion in autonomous vehicles. Now it's selling the unit to Aurora, which makes self-driving tech.







Machine Learning for Energy System

Vast Opportunities

- Load flow analysis/state estimation
- Forecasting (wind, solar, load, prices)
- Fault detection, classification, and localization
- Accelerated market clearing
- Nonlinear control design/RL
- Parameter estimation/Stability assessment
- Many, many, more!

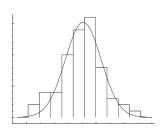
Possible Concerns

Power systems have little room for trial and error! Especially, at fast time scales

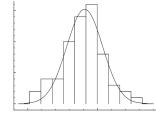


Core challenge: The curse of dimensionality

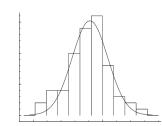
lacktriangle Statistical: Sampling in d dimension with resolution ϵ











Sample complexity:

$$O(\varepsilon^{-d})$$

For $\epsilon=0.1$ and d=100, we would need 10^{100} points.

Computational: Verifying non-negativity of polvnomials

Copositive matrices:

$$[x_1^2 \dots x_d^2] A [x_1^2 \dots x_d^2]^{\mathrm{T}} \ge 0$$

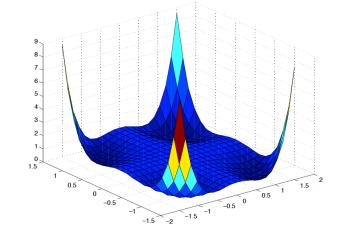
Murty&Kadabi [1987]: Testing co-positivity is NP-Hard

Sum of Squares (SoS):

$$z(x)^T Q z(x) \ge 0$$
, $z_i(x) \in \mathbb{R}[x]$, $x \in \mathbb{R}^d$, $Q \ge 0$

Artin [1927] (Hilbert's 17th problem):

Non-negative polynomials are sum of square of rational functions



Motzkin [1967]:

$$p = x^4y^2 + x^2y^4 + 1 - 3x^2y^2$$

is nonnegative,

not a sum of squares,

but
$$(x^2 + y^2)^2 p$$
 is SoS

Question: Are we asking too much?

• Learnability requires uniform approximation errors across the entire domain

Q: Can we provide local guarantees, and progressively expand as needed?

[arXiv '22] Shen, Bichuch, M

 Lyapunov functions and control barrier functions require strict and exhaustive notions of *invariance*

Q: Can we substitute invariance with less restrictive properties?

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Control synthesis usually aims for the best (optimal) controller

Q: Can we focus on feasibility, rather than optimality?

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[arXiv '21, L4DC 22] Castellano, Min, Bazerque, M

[arXiv 22] Shen, Bichuch, M, Model-free Learning of Regions of Attraction via Recurrent Sets, submitted to CDC 2022, preprint arXiv:2204.10372.

[L4DC 22] Castellano, Min, Bazerque, M, Reinforcement Learning with Almost Sure Constraints, Learning for Dynamics and Control (L4DC) Conference, 2022

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Enrique Mallada (JHU)

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[Submitted on 21 Apr 2022]

Model-free Learning of Regions of Attraction via Recurrent Sets

Yue Shen, Maxim Bichuch, Enrique Mallada





Yue Shen





Maxim Bichuch



Motivation: Estimation of regions of attraction

Having an approximation of the region of attraction allows us to

• Test the limits of controller designs especially for those based on (possibly linear) approximations of nonlinear systems



Verify safety of certain operating condition

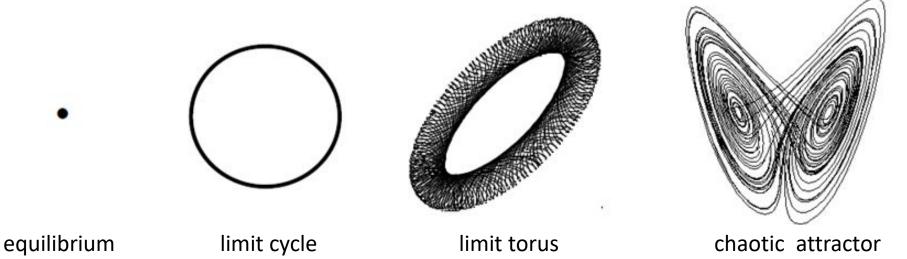


Continuous time dynamical system: $\dot{x}(t) = f(x(t))$

• Initial condition $x_0 = x(0)$, solution at time t: $\phi(t, x_0)$.

Ω-Limit Set Ω(f):
$$x \in \Omega(f) \iff \exists x_0, \{t_n\}_{n \geq 0}, \text{ s.t. } \lim_{n \to \infty} t_n = \infty \text{ and } \lim_{n \to \infty} \phi(t_n, x_0) = x$$

Types of Ω -limit set

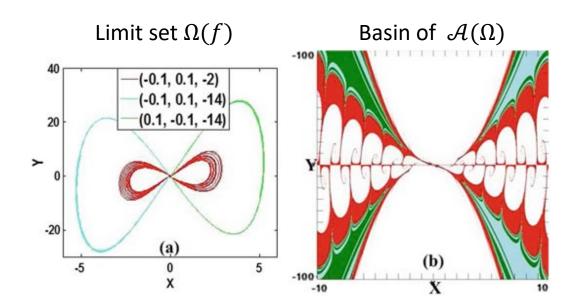


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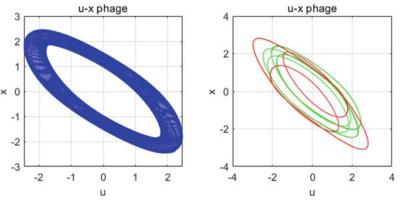
- Initial condition $x_0 = x(0)$, solution at time t: $\phi(t, x_0)$.
- The ω -limit set of the system: $\Omega(f)$

Region of attraction (ROA) of a set $S \subseteq \Omega(f)$:

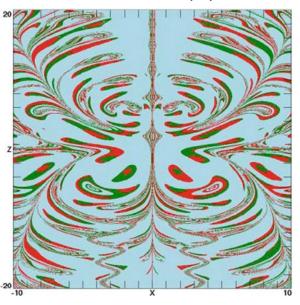
$$\mathcal{A}(S) := \left\{ x_0 \in \mathbb{R}^d | \lim_{t \to \infty} \phi(t, x_0) \in S \right\}$$



Limit set $\Omega(f)$



Basin of $\mathcal{A}(\Omega)$



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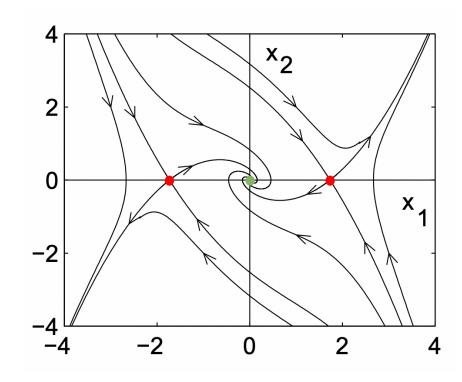
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Simpler Example

$$\begin{bmatrix} \dot{x}_1 \\ \dot{x}_2 \end{bmatrix} = \begin{bmatrix} x_2 \\ -x_1 + \frac{1}{3}x_1^3 - x_2 \end{bmatrix}$$

$$\Omega(f) = \{(0,0), (-\sqrt{3},0), (\sqrt{3},0)\}$$



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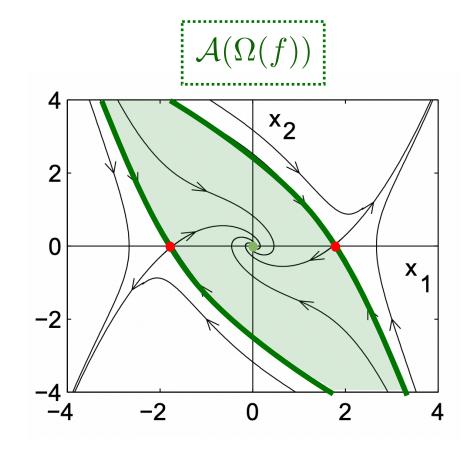
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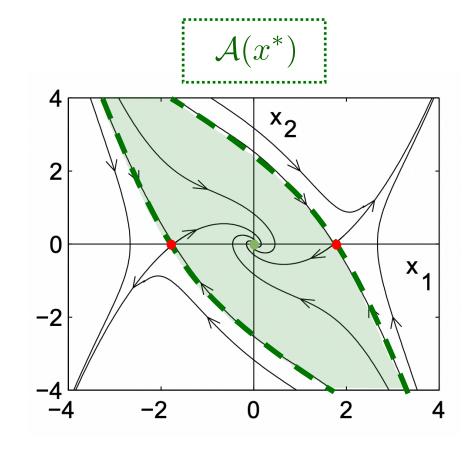
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Asymptotically stable equilibrium at $x^* = (0,0)$



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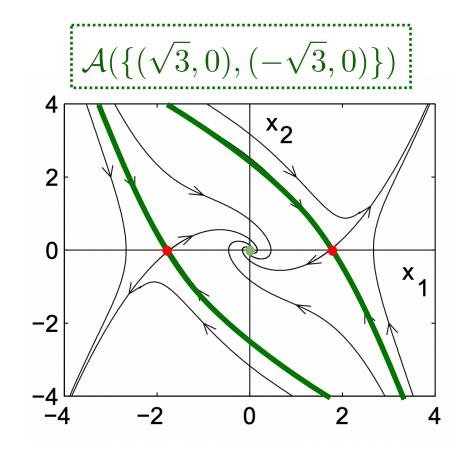
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$$\Omega(f) = \{(0,0), (-\sqrt{3},0), (\sqrt{3},0)\}$$

Unstable equilibria $\{(\sqrt{3},0),(-\sqrt{3},0)\}$



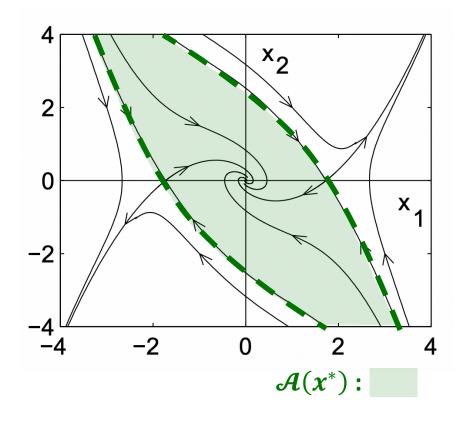
Region of attraction of stable equilibria

Region of attraction (ROA) of a set $S \subseteq \Omega(f)$:

$$\mathcal{A}(S) := \left\{ x_0 \in \mathbb{R}^d | \lim_{t \to \infty} \phi(t, x_0) \in S \right\}$$

Assumption 1. The system $\dot{x}(t) = f(x(t))$ has an asymptotically stable equilibrium at x^* .

Remark 1. It follows from Assumption 1 that the **positively invariant** ROA $\mathcal{A}(x^*)$ is an open contractible **set** [Sontag, 2013], i.e., the identity map of $\mathcal{A}(x^*)$ to itself is null-homotopic [Munkres, 2000].



E. Sontag. "Mathematical Control Theory: Deterministic Finite Dimensional Systems." Springer 2013

J. R. Munkres. "Topology." Prentice Hall 2000

Invariant sets

A set $I \subseteq \mathbb{R}^d$ is **positively invariant** if and only if: $x_0 \in \mathcal{I} \implies \phi(t, x_0) \in \mathcal{I}, \quad \forall t \in \mathbb{R}^+$ Any trajectory starting in the set remains in inside it

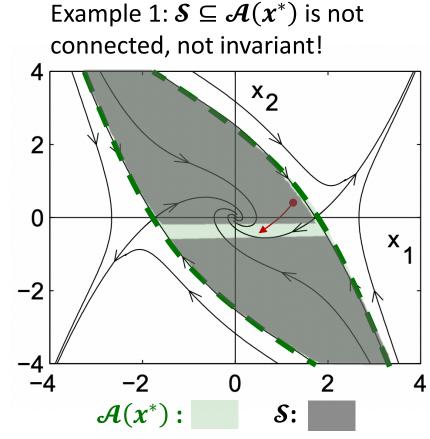
- Invariant sets guarantee stability
 Lyapunov stability: solutions starting "close enough" to the equilibrium (within a distance δ) remain "close enough" forever (within a distance ε))
- Invariant sets further certify asymptotic stability via Lyapunov's direct method Asymptotic stability: solutions that start close enough not only remain close enough but also eventually converge to the equilibrium.)

 Regions of attraction are invariant sets, and so are the outcome of most approximation methods!

Challenges of working with invariant set

Learning ROA $\mathcal{A}(x^*)$ by finding an invariant set $\mathcal{S} \subseteq \mathcal{A}(x^*)$

- **S** is topologically constrained
 - If $S \cap \Omega(f) = \{x^*\}$, then S is connected

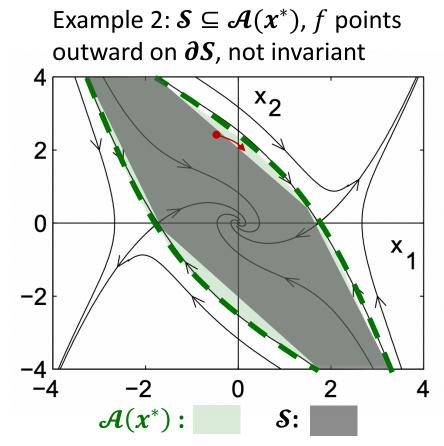


A not invariant trajectory: • • •

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- **S** is geometrically constrained
 - f should point inwards for $x \in \partial S$

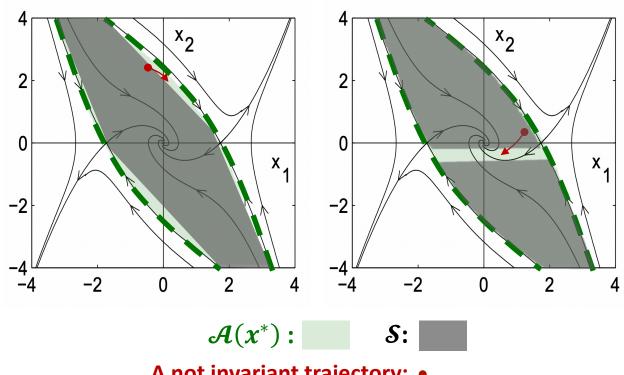


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A subset of an invariant set is not necessary an invariant set



A not invariant trajectory: •

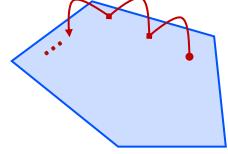
Recurrent sets: Letting things go, and come back

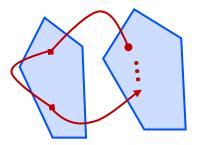
A set $\mathcal{R} \subseteq \mathbb{R}^d$ is **recurrent** if for any $x_0 \in \mathcal{R}$, whenever $\phi(t, x_0) \notin \mathcal{R}$, $t \ge 0$, then $\exists t' > t$ such that $\phi(t', x_0) \in \mathcal{R}$.

Property of Recurrent Sets

- R need not be connected
- $\mathcal R$ does **not** require f to **point inwards** on all $\partial \mathcal R$

Recurrent sets, while not invariant, guarantee that solutions that start in this set, will come back **infinitely often, forever!**





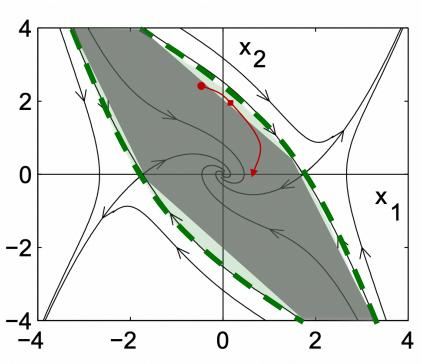
Recurrent set \mathcal{R} :

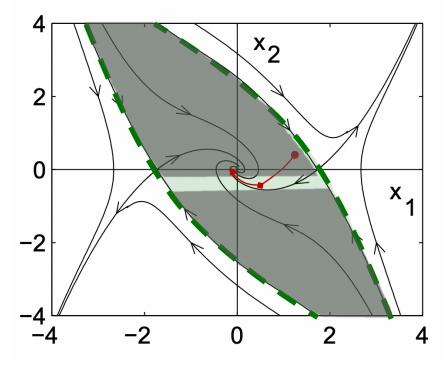
A recurrent trajectory:

Recurrent sets: Letting things go, and come back

A set $\mathcal{R} \subseteq \mathbb{R}^d$ is **recurrent** if for any $x_0 \in \mathcal{R}$, whenever $\phi(t, x_0) \notin \mathcal{R}$, $t \ge 0$, then $\exists t' > t$ such that $\phi(t', x_0) \in \mathcal{R}$.

Previous two good inner approximations of $\mathcal{A}(x^*)$ are recurrent sets

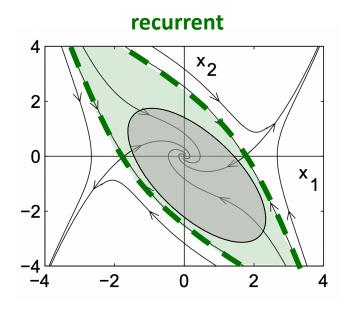


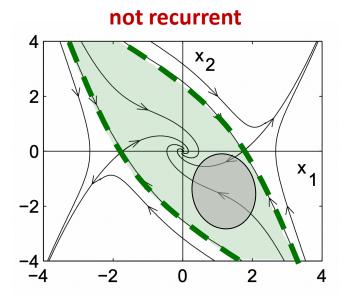


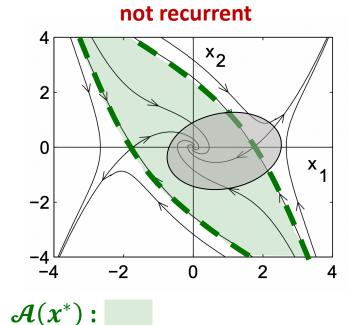
A set $\mathcal{R} \subseteq \mathbb{R}^d$ is **recurrent** if for $x_0 \in \mathcal{R}$, $\phi(t, x_0) \notin \mathcal{R} \Rightarrow \exists t' > t$, s.t. $\phi(t', x_0) \in \mathcal{R}$

Theorem 2. Let $\mathcal{R} \subset \mathbb{R}^d$ be a <u>compact</u> set satisfying $\partial \mathcal{R} \cap \Omega(f) = \emptyset$. Then:

$$\mathcal{R}$$
 is recurrent $\longleftrightarrow \begin{array}{c} \mathcal{R} \cap \Omega(f) \neq \emptyset \\ \mathcal{R} \subset \mathcal{A}(\mathcal{R} \cap \Omega(f)) \end{array}$





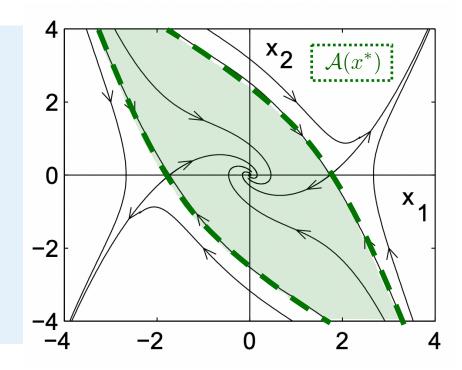


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Assumption 2. The ω -limit set $\Omega(f)$ is composed by **hyperbolic equilibrium points**, with only one of them, say x^* , being asymptotically stable.

Corollary 2. Let Assumptions 1 and 2 hold, and $\mathcal{R} \subset \mathbb{R}^d$ be a <u>compact</u> set satisfying $\partial \mathcal{R} \cap \Omega(f) = \emptyset$. Then:

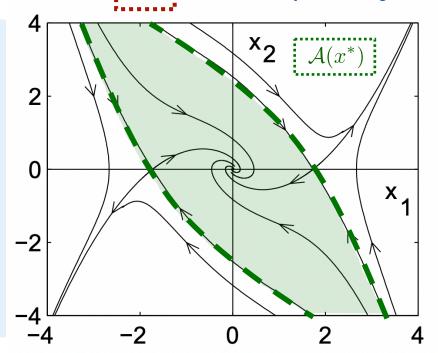
 \mathcal{R} is recurrent \iff $\mathcal{R} \cap \Omega(f) = \{x^*\}$ $\mathcal{R} \subset \mathcal{A}(x^*)$



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$$\mathcal{R} \text{ is recurrent} \longleftrightarrow \begin{array}{c} \mathcal{R} \cap \Omega(f) = \{x^*\} \\ \mathcal{R} \subset \mathcal{A}(x^*) \end{array}$$



Idea: Use recurrence as a mechanism for finding inner approximations of $\mathcal{A}(x^*)$

Potential Issues:

- We do not know how long it takes to come back!
- We need to adapt results to trajectory samples

au-recurrent sets

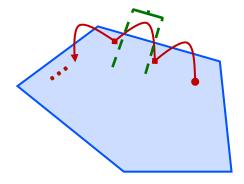
Time elapsed ≤ τ

A set \mathcal{R} is τ -recurrent if whenever $x_0 \in \mathcal{R}$, $\exists t' \in (0, \tau]$ s.t. $\phi(t', x_0) \in \mathcal{R}$

Theorem 3. Under Assumption 1, any compact set \mathcal{R} satisfying:

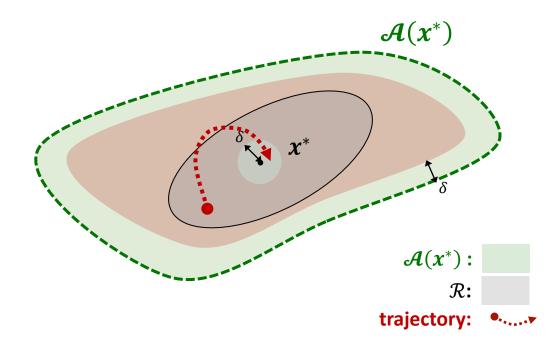
$$x^* + \mathcal{B}_{\delta} \subseteq \mathcal{R} \subseteq \mathcal{A}(x^*) \setminus \{\partial \mathcal{A}(x^*) + \text{int } \mathcal{B}_{\delta}\}$$

is τ -recurrent for $\tau \geq \bar{\tau}(\delta) \coloneqq \frac{\underline{c}(\delta) - \bar{c}(\delta)}{a(\delta)}$.



 τ -recurrent set \mathcal{R} :

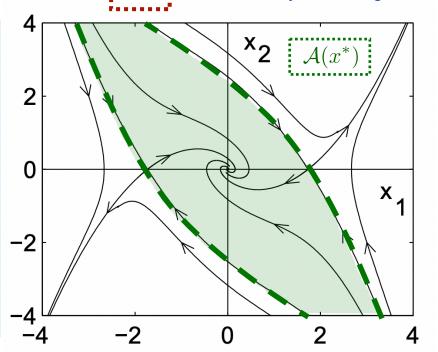
trajectory:



A set $\mathcal{R} \subseteq \mathbb{R}^d$ is **recurrent** if for $x_0 \in \mathcal{R}$, $\phi(t, x_0) \notin \mathcal{R} \Rightarrow \exists t' > t$, s.t. $\phi(t', x_0) \in \mathcal{R}$

Corollary 2. Let Assumptions 1 and 2 hold, and $\mathcal{R} \subset \mathbb{R}^d$ be a <u>compact</u> set satisfying $\partial \mathcal{R} \cap \Omega(f) = \emptyset$. Then:

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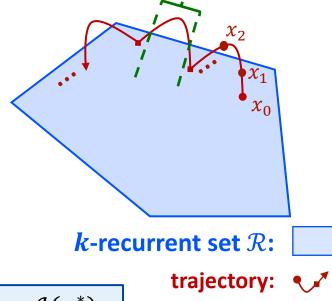
• We need to adapt results to trajectory samples

Learning recurrent sets from k-length trajectory samples

Consider finite length trajectories:

$$x_n = \phi(n\tau_s, x_0), \quad x_0 \in \mathbb{R}^d, n \in \mathbb{N},$$
 where $\tau_s > 0$ is the sampling period.

• A set $\mathcal{R} \subseteq \mathbb{R}^d$ is k-recurrent if whenever $x_0 \in \mathcal{R}$, then $\exists n \in \{1, ..., k\}$ s.t. $x_n \in \mathcal{R}$



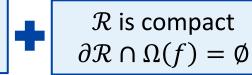
steps elapsed $\leq k$

(time elapsed $\leq k\tau_s$)

Sufficiency:

$$\mathcal{R}$$
 is k -recurrent

$$\mathcal{R}$$
 is au -recurrent with $au=k au_s$



 $\Rightarrow \qquad \mathcal{R} \subset \mathcal{A}(x^*)$

(Corollary 2, under Assumption 2)

Necessity:

Theorem 4. Under Assumption 1, any compact set \mathcal{R} satisfying:

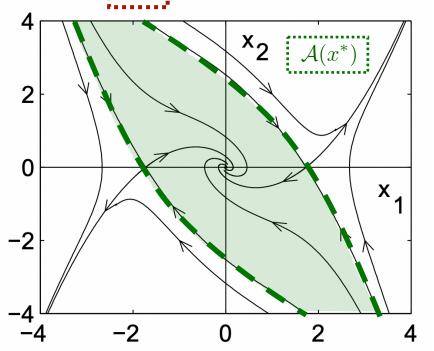
$$\mathcal{B}_{\delta} + x^* \subseteq \mathcal{R} \subseteq \mathcal{A}(x^*) \setminus \{\partial \mathcal{A}(x^*) + \text{int } \mathcal{B}_{\delta}\}$$

is k-recurrent for any $k > \bar{k} := \bar{\tau}(\delta)/\tau_s$.

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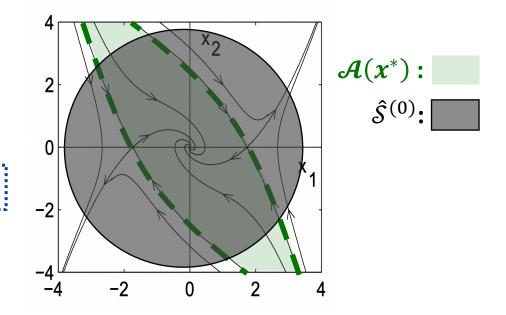
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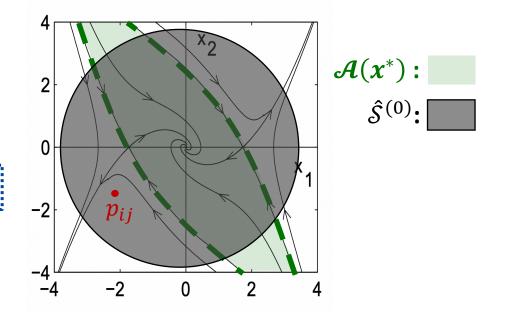
Algorithm:

• Initialize $\hat{\mathcal{S}}^{(0)}$ as $\hat{\mathcal{S}}^{(0)}\coloneqq\{x|\|x\|_2\leq b^{(0)}\coloneqq c\}\supseteq\mathcal{B}_\delta$



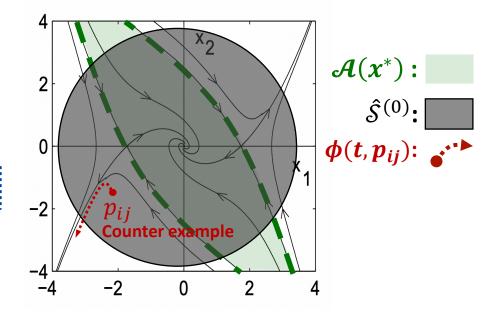
Algorithm:

- Initialize $\hat{\mathcal{S}}^{(0)}$ as $\hat{\mathcal{S}}^{(0)}\coloneqq\{x|\|x\|_2\leq b^{(0)}\coloneqq c\}\supseteq\mathcal{B}_\delta$
- For iteration i = 0,1,... do: (set updates)
 - For iteration j = 0,1, ... do: (samples)
 - Generate random sample $p_{ij} \in \hat{\mathcal{S}}^{(i)}$ uniformly



Algorithm:

- Initialize $\hat{\mathcal{S}}^{(0)}$ as $\hat{\mathcal{S}}^{(0)}\coloneqq\{x|\|x\|_2\leq b^{(0)}\coloneqq c\}\supseteq\mathcal{B}_\delta$
- **For** *iteration* i = 0,1,... **do:**
 - *For* iteration j = 0,1, ... do:
 - Generate random sample $p_{ij} \in \hat{\mathcal{S}}^{(i)}$ uniformly
 - If p_{ij} is a counter-example w.r.t $\hat{S}^{(i)}$ do:

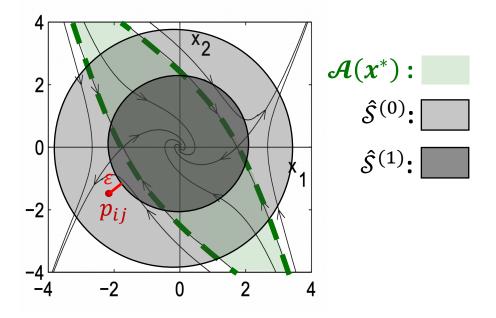


We say sample point p_{ij} is a valid k-recurrent point w.r.t current approximation $\hat{\mathcal{S}}^{(i)}$ if starting from $x_0 = p_{ij}$, $\exists \ n \in \{1, \dots, k\}, \ \text{s.t.} \ x_n \in \hat{\mathcal{S}}^{(i)}.$

Otherwise, we say p_{ij} is a counter-example.

Algorithm:

- Initialize $\hat{\mathcal{S}}^{(0)}$ as $\hat{\mathcal{S}}^{(0)}\coloneqq\{x|\|x\|_2\leq b^{(0)}\coloneqq c\}\supseteq\mathcal{B}_{\delta}$
- **For** *iteration* i = 0,1, ... **do:**
 - *For* iteration j = 0,1, ... do:
 - Generate random sample $p_{ij} \in \hat{\mathcal{S}}^{(i)}$ uniformly
 - If p_{ij} is a counter-example w.r.t $\hat{S}^{(i)}$ do:
 - *Update* $b^{(i)}$ *to* $b^{(i+1)}$, $\hat{S}^{(i)}$ *to* $\hat{S}^{(i+1)}$



We say sample point p_{ij} is a valid k-recurrent point w.r.t current approximation $\hat{\mathcal{S}}^{(i)}$ if starting from $x_0 = p_{ij}$, $\exists n \in \{1, ..., k\}$, s.t. $x_n \in \hat{\mathcal{S}}^{(i)}$.

Otherwise, we say p_{ij} is a counter-example.

If p_{ij} is a counter-example, we update:

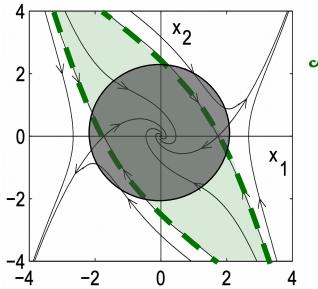
$$b^{(i+1)} = \|p_{ij}\|_{2} - \varepsilon;$$

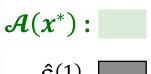
$$\hat{S}^{(i+1)} = \{x | \|x\|_{2} \le b^{(i+1)} \},$$

where $\varepsilon > 0$ is an algorithm parameter expressing the level of conservativeness in our update.

Algorithm:

- Initialize $\hat{\mathcal{S}}^{(0)}$ as $\hat{\mathcal{S}}^{(0)}\coloneqq\{x|\|x\|_2\leq b^{(0)}\coloneqq c\}\supseteq\mathcal{B}_\delta$
- **For** *iteration* i = 0,1,... **do:**
 - *For* iteration j = 0,1, ... do:
 - Generate random sample $p_{ij} \in \hat{\mathcal{S}}^{(i)}$ uniformly
 - If p_{ij} is a counter-example w.r.t $\hat{S}^{(i)}$ do:
 - Update $b^{(i)}$ to $b^{(i+1)}$, $\hat{S}^{(i)}$ to $\hat{S}^{(i+1)}$
 - Break
 - End if
 - End for
- End for





We say sample point p_{ij} is a valid k-recurrent point w.r.t current approximation $\hat{\mathcal{S}}^{(i)}$ if starting from $x_0 = p_{ij}$, $\exists n \in \{1, ..., k\}$, s.t. $x_n \in \hat{\mathcal{S}}^{(i)}$.

Otherwise, we say p_{ij} is a counter-example.

If p_{ij} is a counter-example, we update:

$$b^{(i+1)} = \|p_{ij}\|_{2} - \varepsilon;$$

$$\hat{S}^{(i+1)} = \{x | \|x\|_{2} \le b^{(i+1)} \},$$

where $\varepsilon > 0$ is an algorithm parameter expressing the level of conservativeness in our update.

Parameter choice

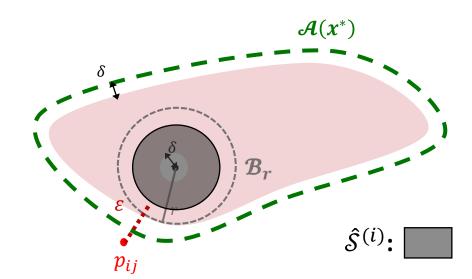
Choice of ε : $b^{(i+1)} = ||p_{ij}|| - \varepsilon$

• Given $k > \overline{k}$, any set $S^{(i)} = \{x : ||x|| \le b^{(i)}\}$ satisfying:

$$\mathcal{B}_{\delta} \subseteq \mathcal{S}^{(i)} \subseteq \mathcal{A}(0) \setminus \{\partial \mathcal{A}(0) + \text{int } \mathcal{B}_{\delta}\}\$$

is k-recurrent.

- Let \mathcal{B}_r the largest ball inside $\mathcal{A}(0)\setminus\{\partial\mathcal{A}(0) + \text{int } \mathcal{B}_\delta\}$
- Then, if $\varepsilon \leq r \delta$ we always guarantee $\mathcal{B}_{\delta} \subseteq \mathcal{S}^{(i)}$



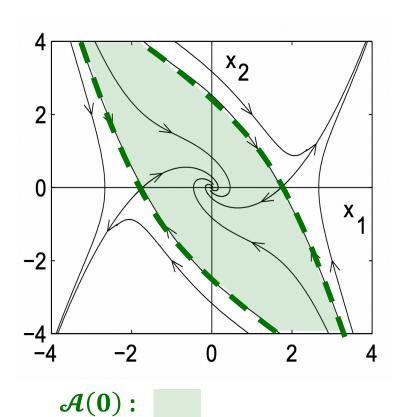
Choice of trajectory length k:

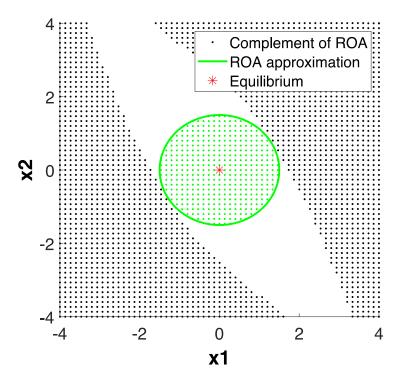
- $\bar{k}(\delta)$ depends highly non-trivially on δ .
- If $k < \overline{k}(\delta)$, we get $b^{(i)} < 0 \Longrightarrow$ Failure!
- Solution: doubling the size of k, i.e., $k^+ = 2k$, every time we fail.

With $oldsymbol{k}$ -doubling, the total number of counter-examples is bounded by

$$\# counter-examples \leq \frac{b^{(0)}}{\varepsilon} \log_2 \overline{k}(\delta)$$

Algorithm Result - Sphere Approximations





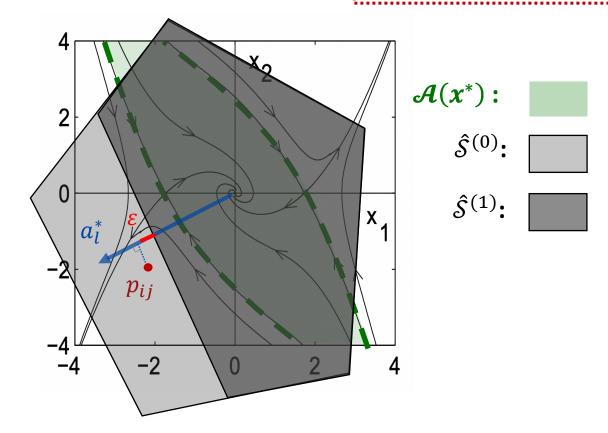
Polytope approximations of RoA

Algorithm:

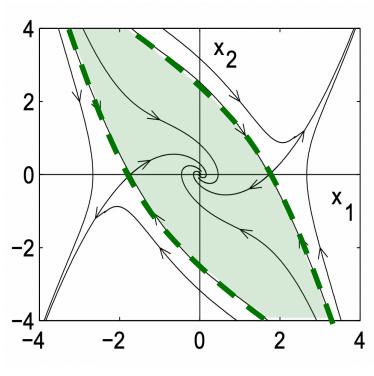
Initialize $\hat{\mathcal{S}}^{(0)}$ as $\hat{\mathcal{S}}^{(0)} \coloneqq \{x | Ax \leq b^{(0)} \coloneqq c \mathbb{I}_n\} \supseteq \mathcal{B}_{\delta}$

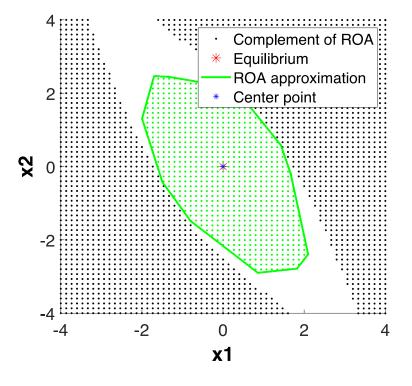
$$\hat{S}^{(0)} \coloneqq \{x | Ax \le b^{(0)} \coloneqq c \mathbb{I}_n\} \supseteq \mathcal{B}_{\delta}$$

Exploration direction matrix $A := [a_1, ..., a_n] \subseteq$ $\mathbb{R}^{n \times d}$, where each row vector a_l is a normalized exploration direction indexed by $l \in$ $\{1, ..., n\}.$



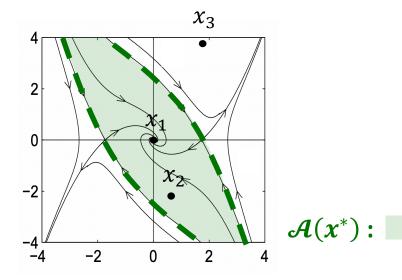
Algorithm Result – Polytope Approximation



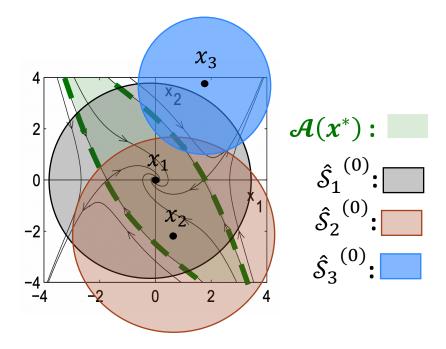


 $\mathcal{A}(\mathbf{0})$:

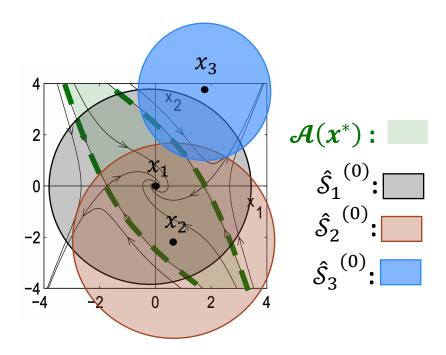
- Consider $h \in \mathbb{N}^+$ center points x_q indexed by $q \in \{1, ..., h\}$.
 - Let the first center point $x_1 = x^* = 0$
 - Additional center point x_2, \dots, x_h can be designed chosen uniformly.



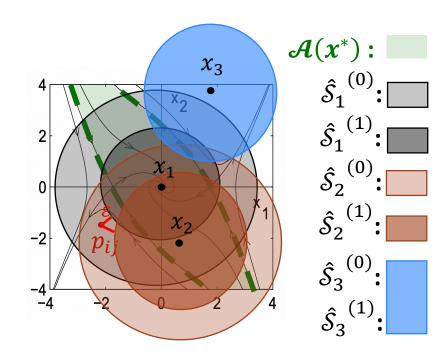
- Consider $h \in \mathbb{N}^+$ center points x_q indexed by $q \in \{1, ..., h\}$.
 - Let the first center point $x_1 = x^* = 0$
 - Additional center point $x_2, ..., x_h$ can be designed chosen uniformly.
- Respectively defined approximations centered at each x_q
 - (Sphere case) $\hat{S}_q^{(i)} := \{x | ||x x_q||_2 \le b_q^{(i)}\}$
 - (Polytope case) $\hat{\mathcal{S}}_q^{(i)} \coloneqq \{x | A(x x_q) \le b_q^{(i)}\}$



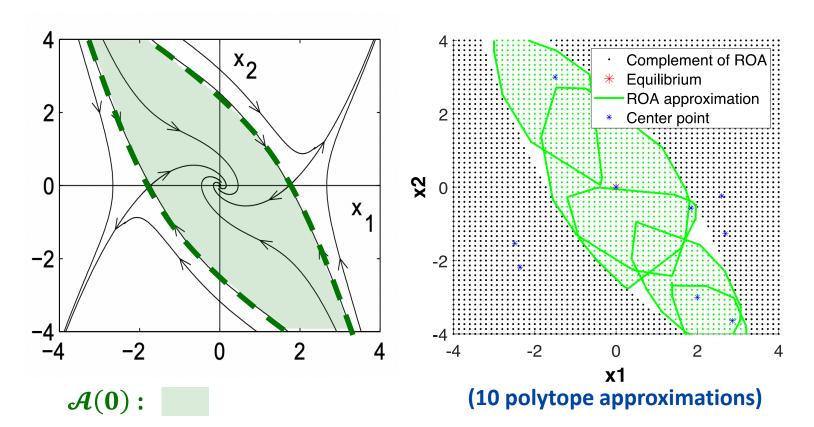
- Consider $h \in \mathbb{N}^+$ center points x_q indexed by $q \in \{1, ..., h\}$.
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 - (Polytope case) $\hat{\mathcal{S}}_q^{(i)} \coloneqq \{x | A(x x_q) \le b_q^{(i)}\}$
- Multiple centers approximation $\hat{\mathcal{S}}_{ ext{multi}}^{(i)} := \cup_{q=1}^h \hat{S}_q^{(i)}$

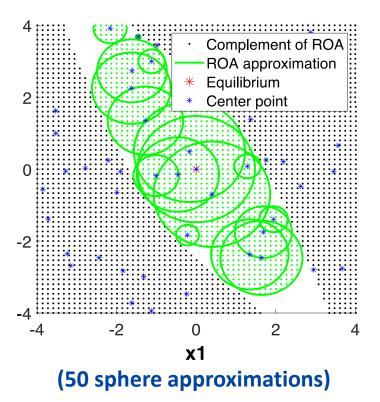


- Consider $h \in \mathbb{N}^+$ center points x_q indexed by $q \in \{1, ..., h\}$.
 - Let the first center point $x_1 = x^* = 0$
 - Additional center point $x_2, ..., x_h$ can be designed chosen uniformly.
- Respectively defined approximations centered at each x_q
 - (Sphere case) $\hat{S}_q^{(i)} := \{x | ||x x_q||_2 \le b_q^{(i)}\}$
 - (Polytope case) $\hat{\mathcal{S}}_q^{(i)} \coloneqq \{x | A(x x_q) \le b_q^{(i)}\}$
- Multiple centers approximation $\,\hat{\mathcal{S}}^{(i)}_{
 m multi} := \cup_{q=1}^h \hat{S}^{(i)}_q \,$
- If $\mathbf{p_{ij}}$ is a counter-example w.r.t $\hat{\mathcal{S}}_{\mathrm{multi}}^{(i)}$
 - We shrink every $\hat{\mathcal{S}}_q^{\;(i)}$ satisfying $p_{ij} \in \hat{\mathcal{S}}_q^{\;(i)}$
 - For the rest approximations, we simply let $\hat{\mathcal{S}}_q^{~(i+1)} = \hat{\mathcal{S}}_q^{~(i)}$

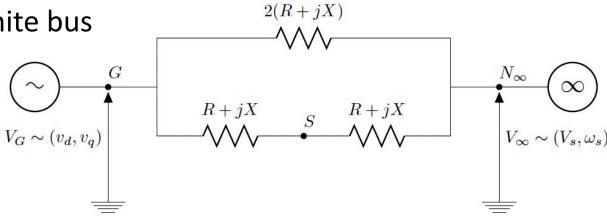


Algorithm results – Multi-center approximation

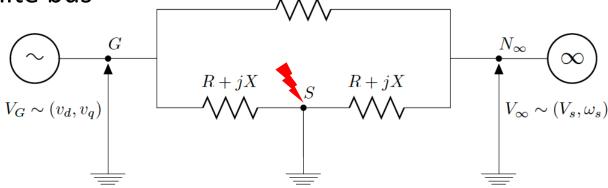




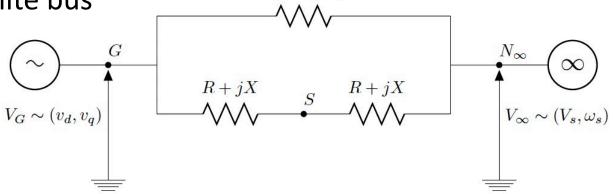
• Synchronous machine connected to infinite bus



- Synchronous machine connected to infinite bus
- t_1 lower line is short-circuited



- Synchronous machine connected to infinite bus
- t_1 lower line is short-circuited
- t₂ fault is cleared



- Synchronous machine connected to infinite bus
- t_1 lower line is short-circuited
- t₂ fault is cleared

$$\frac{d\delta}{dt} = \omega - \omega_s$$

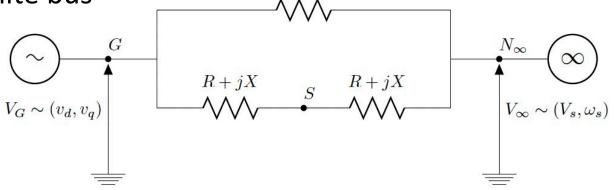
$$2H\frac{d\omega}{dt} = P_m - (v_d i_d + v_q i_q + e i_d^2 + r i_q^2)$$

$$T'_{d_0} \frac{de'_q}{dt} = -e'_q - (x_d - x'_d) i_d + E_{fd}$$

$$T_a \frac{dE_{fd}}{dt} = -E_{fd} + K_a (V_{ref} - V_t)$$

$$T_g \frac{dP_m}{dt} = -P_m + P_{ref} + K_g (\omega_{ref} - \omega)$$

$$i_q = \frac{(X - x'_d) V_s \sin(\delta) - (R + r) (V_s \cos(\delta) - e'_q)}{(R + r)^2 + (X + x'_d) (X + x_q)}$$



$$i_d = \frac{X - x_q}{R + r} i_q - \frac{1}{R + r} V_s \sin(\delta)$$

$$v_d = x_q i_q - r - i_d$$

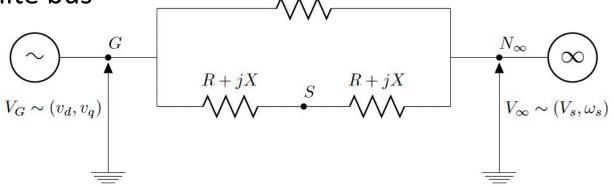
$$v_q = R i_q + X i_d + V_s \cos(\delta)$$

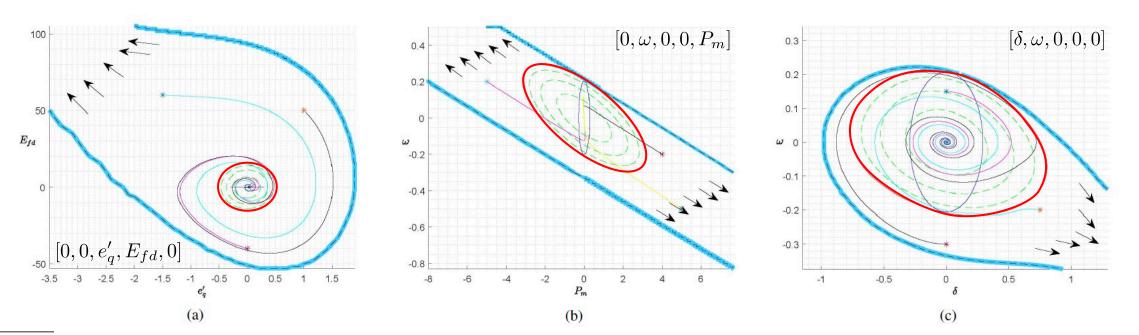
$$V_t = \sqrt{v_d^2 + v_q^2}$$

$$T'_{d_0} = 9.67$$
 $x_d = 2.38$ $x'_d = 0.336$ $x_q = 1.21$ $H = 3$ $r = 0.002$ $\omega_s = \omega_{ref} = 1$ $R = 0.01$ $X = 1.185$ $V_s = 1$ $T_a = 1$ $K_a = 70$ $V_{ref} = 1$ $T_g = 0.4$ $K_g = 0.5$ $P_{ref} = 0.7$

- Synchronous machine connected to infinite bus
- t_1 lower line is short-circuited
- t₂ fault is cleared

SoS approx. in red (2d-sections)



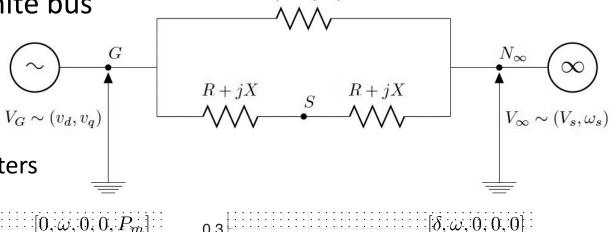


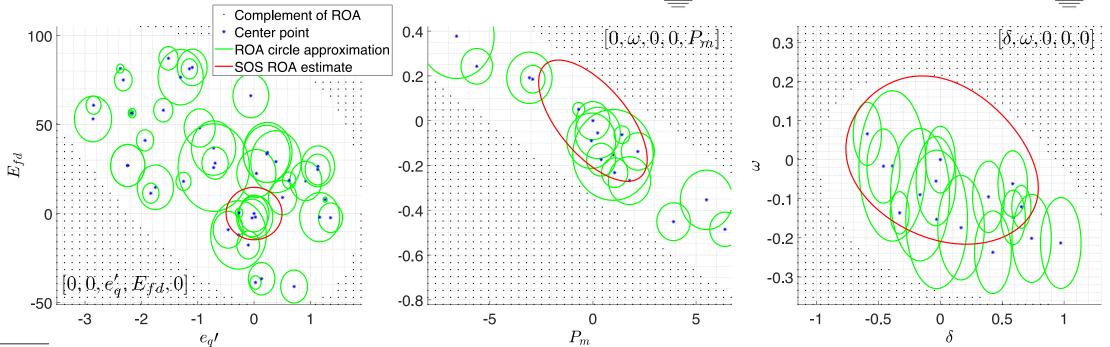
M. Tacchi et al "Power system transient stability analysis using SoS programming" Power System Computation Conference (PSCC) 2018

• Synchronous machine connected to infinite bus

- t_1 lower line is short-circuited
- t₂ fault is cleared

Multi-center in green: $\tau_s = 1$, k = 40, 2.5K centers



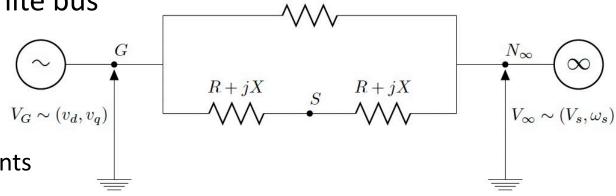


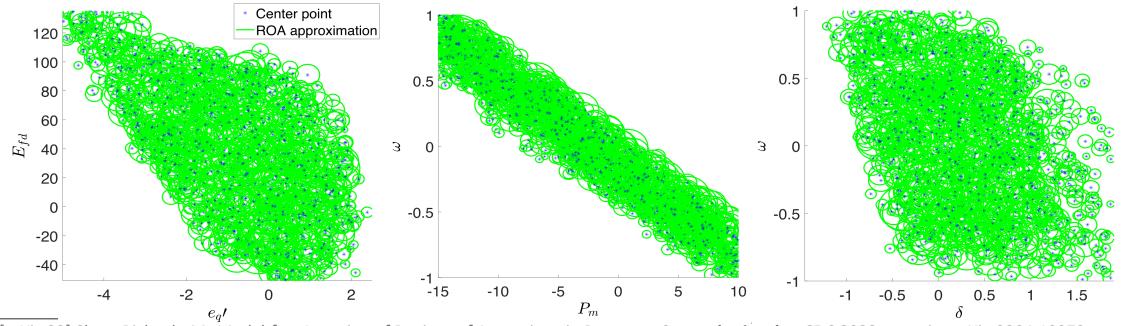
[arXiv 22] Shen, Bichuch, M, Model-free Learning of Regions of Attraction via Recurrent Sets, submitted to CDC 2022, preprint arXiv:2204.10372.

Synchronous machine connected to infinite bus

- t_1 lower line is short-circuited
- t₂ fault is cleared

Multi-center in green: $\tau_s = 1$, k = 40, 1.5K points





[arXiv 22] Shen, Bichuch, M, Model-free Learning of Regions of Attraction via Recurrent Sets, submitted to CDC 2022, preprint arXiv:2204.10372.

Conclusions and Future work

Take-aways

- Proposed a relaxed notion of invariance known as recurrence.
- Provide necessary and sufficient conditions for a recurrent set to be an innerapproximation of the ROA.
- Our algorithms are sequential, and only incur a limited number of counter-examples.

Ongoing work

Sample complexity bounds, smart choice of multi-points, control recurrent sets

Thanks!

Related Publication:

[arXiv 22] Shen, Bichuch, M, Model-free Learning of Regions of Attraction via Recurrent Sets, submitted to CDC 2022, preprint arXiv:2204.10372.



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